



Lecture #15

OUTLINE

- Digital Representation
- Logic
- Binary representations
- Combinatorial logic circuits

Reading

- Hambley 7-7.5
- Rabaey 5.1-5.3.2, 5.4-5.4.2, 5.5-5.5.2



Digital Circuits – Introduction

- Analog: signal amplitude is continuous with time.
- Digital: signal amplitude is represented by a restricted set of discrete numbers.
 - Binary: only two values are allowed to represent the signal: High or low (i.e. logic 1 or 0).
- Digital word:
 - Each binary digit is called a bit
 - A series of bits form a word
 - Byte is a word consisting of 8-bits
- Advantages of digital signal
 - Digital signal is more resilient to noise → can more easily differentiate high (1) and low (0)
- Transmission
 - Parallel transmission over a bus containing n wires.
 - Faster but short distance (internal to a computer or chip)
 - Serial transmission (transmit bits sequentially)
 - Longer distance



Digital Signal Representations

Binary numbers can be used to represent any quantity.

We generally have to agree on some sort of “code”, and the resolution of the signal in order to know the form and the number of binary digits (“bits”) required.

Example 1: Voltage signal with maximum value 2 Volts

- Binary two (**10**) could represent a 2 Volt signal.
- To encode the signal to an accuracy of 1 part in 64 (1.5% precision), 6 binary digits (“bits”) are needed

Example 2: Sine wave signal of known frequency and maximum amplitude 50 μV ; 1 μV resolution needed.



Why Digital?

(For example, why CDROM audio vs. vinyl recordings?)

- Digital signals can be transmitted, received, amplified, and re-transmitted with no degradation.
- Digital information is easily and inexpensively stored (in RAM, ROM, *etc.*), with arbitrary accuracy.
- Complex logical functions are easily expressed as binary functions (*e.g.* in control applications).
- Digital signals are easy to manipulate (as we shall see).



Reminder About Binary and Decimal Numbering Systems

$$\begin{aligned} 110001_2 &= 1 \times 2^5 + 1 \times 2^4 + 0 \times 2^3 + 0 \times 2^2 + 0 \times 2^1 + 1 \times 2^0 \\ &= 32_{10} + 16_{10} + 1_{10} \\ &= 49_{10} \\ &= 4 \times 10^1 + 9 \times 10^0 \end{aligned}$$



Example 2 (continued)

Possible digital representation for the sine wave signal:

Analog representation: Amplitude in μV	Digital representation: Binary number
1	000001
2	000010
3	000011
4	000100
5	000101
8	001000
16	010000
32	100000
50	110010
63	111111



Binary Representation

- N bit can represent 2^N values: typically from 0 to $2^N - 1$
 - 3-bit word can represent 8 values: e.g. 0, 1, 2, 3, 4, 5, 6, 7
- Conversion
 - Integer to binary
 - Fraction to binary ($13.5_{10} = 1101.1_2$ and $0.392_{10} = 0.011001_2$)
- Octal and hexadecimal



Digital Representations of Logical Functions

Digital signals offer an easy way to perform logical functions, using Boolean algebra.

- Variables have two possible values: “**true**” or “**false**”
 - usually represented by **1** and **0**, respectively.

All modern control systems use this approach.

Example: Hot tub controller with the following algorithm

Turn on the heater if the temperature is less than desired ($T < T_{\text{set}}$) **and** the motor is on **and** the key switch to activate the hot tub is closed.



Notation for Logical Expressions

Basic logical functions:

AND:	"dot"	Example: $X = A \cdot B$
OR:	"+" sign	Example: $Y = \overline{A} + B$
NOT:	"bar over symbol"	Example: $Z = \overline{A}$

- **Any logical expression can be constructed using these basic logical functions**

Additional logical functions:

Inverted AND = NAND:	\overline{AB}	(only 0 when A and $B=1$)
Inverted OR = NOR:	$\overline{A + B}$	(only 1 when $A = B = 0$)
Exclusive OR:	$A \oplus B$	(only 1 when A, B differ) i.e., $A + B$ except $A \cdot B$



Boolean algebras

- are algebraic structures which "capture the essence" of the logical operations AND, OR and NOT as well as the corresponding set theoretic operations intersection, union and complement.
- They are named after George Boole, an English mathematician at University College Cork, who first defined them as part of a system of logic in the mid 19th century. Specifically, Boolean algebra was an attempt to use algebraic techniques to deal with expressions in the propositional calculus. Today, Boolean algebras find many applications in electronic design. They were first applied to switching by Claude Shannon in the 20th century.



Boolean algebras

- The operators of Boolean algebra may be represented in various ways. Often they are simply written as AND, OR and NOT.
- In describing circuits, NAND (NOT AND), NOR (NOT OR) and XOR (eXclusive OR) may also be used.
- Mathematicians often use + for OR and \cdot for AND (since in some ways those operations are analogous to addition and multiplication in other algebraic structures) and represent NOT by a line drawn above the expression being negated.

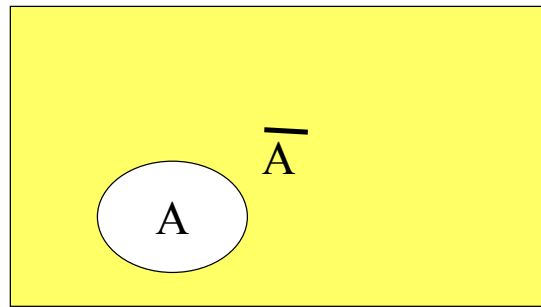


Boolean Algebra

- NOT operation (inverter) $A \bar{A} = 0$
 $A + \bar{A} = 1$
- AND operation $A \bar{A} = A$
 $A \bar{1} = A$
 $A \bar{0} = 0$
 $A \bar{B} = B \bar{A}$
 $(A \bar{B}) \bar{C} = A \bar{(B \bar{C})}$
- OR operation $A + A = A$
 $A + 1 = 1$
 $A + 0 = A$
 $A + B = B + A$
 $(A + B) + C = A + (B + C)$



Graphic Representation



$$A \cap \bar{A} = 0$$

$$A + \bar{A} = 1$$

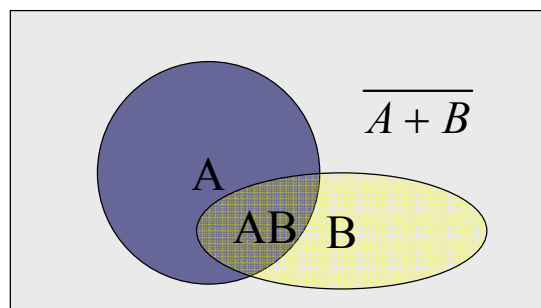
Full square = complete set = 1

Yellow part = NOT(A) = \bar{A}

White circle = A



Graphic Representation



$$A \oplus B = A\bar{B} + \bar{A}B = (A + B) \cap (\bar{A} + \bar{B}) = \overline{A \cap B} = \overline{AB}$$

Exclusive OR = yellow and blue part –

intersection/overlap part

= exactly when only one of the input is true



Boolean Algebra

- Distributive Property

$$A \cdot (B + C) = A \cdot B + A \cdot C$$

$$(A + B) \cdot C = (A + B) \cdot (A + C)$$

- De Morgan's laws

$$\overline{A + B} = \overline{A} \cdot \overline{B}$$

$$\overline{A \cdot B} = \overline{A} + \overline{B}$$

- An excellent web site to visit

□ http://en.wikipedia.org/wiki/Boolean_algebra



Examples

$$F = A \cdot \overline{B} \cdot C + A \cdot B \cdot C + (C + D) \cdot (\overline{D} + E)$$

$$F = C \cdot (A + \overline{D} + E) + D \cdot E$$



Logic Gates and Memories

■ Logic gates

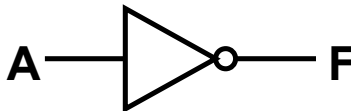
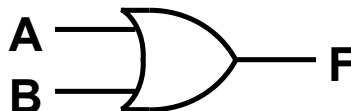
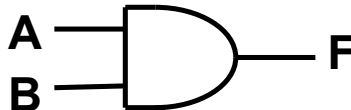
- Combine several logic variable inputs to produce a logic variable output

■ Memory

- Memoryless: output at a given instant depends the input values of that instant.
- Memory: output depends on previous and present input values.



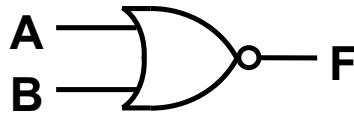
Logic Functions, Symbols, & Notation

<u>NAME</u>	<u>SYMBOL</u>	<u>NOTATION</u>	<u>TRUTH TABLE</u>															
“NOT”		$F = \overline{A}$	<table><tr><th>A</th><th>F</th></tr><tr><td>0</td><td>1</td></tr><tr><td>1</td><td>0</td></tr></table>	A	F	0	1	1	0									
A	F																	
0	1																	
1	0																	
“OR”		$F = A+B$	<table><tr><th>A</th><th>B</th><th>F</th></tr><tr><td>0</td><td>0</td><td>0</td></tr><tr><td>0</td><td>1</td><td>1</td></tr><tr><td>1</td><td>0</td><td>1</td></tr><tr><td>1</td><td>1</td><td>1</td></tr></table>	A	B	F	0	0	0	0	1	1	1	0	1	1	1	1
A	B	F																
0	0	0																
0	1	1																
1	0	1																
1	1	1																
“AND”		$F = A \cdot B$	<table><tr><th>A</th><th>B</th><th>F</th></tr><tr><td>0</td><td>0</td><td>0</td></tr><tr><td>0</td><td>1</td><td>0</td></tr><tr><td>1</td><td>0</td><td>0</td></tr><tr><td>1</td><td>1</td><td>1</td></tr></table>	A	B	F	0	0	0	0	1	0	1	0	0	1	1	1
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Logic Functions, Symbols, & Notation 2

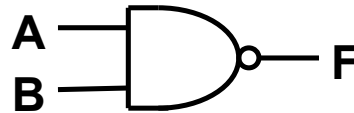
“NOR”



$$F = \overline{A+B}$$

A	B	F
0	0	1
0	1	0
1	0	0
1	1	0

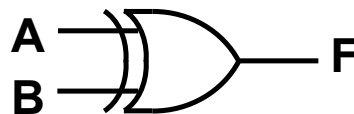
“NAND”



$$F = \overline{A \cdot B}$$

A	B	F
0	0	1
0	1	1
1	0	1
1	1	0

“XOR”
(exclusive OR)



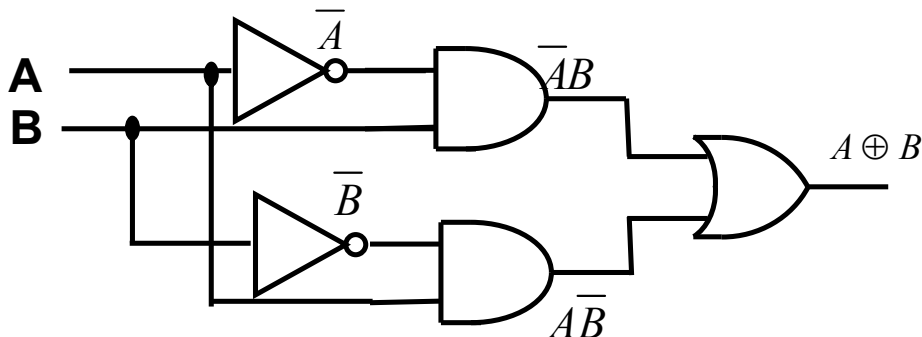
$$F = A \oplus B$$

A	B	F
0	0	0
0	1	1
1	0	1
1	1	0



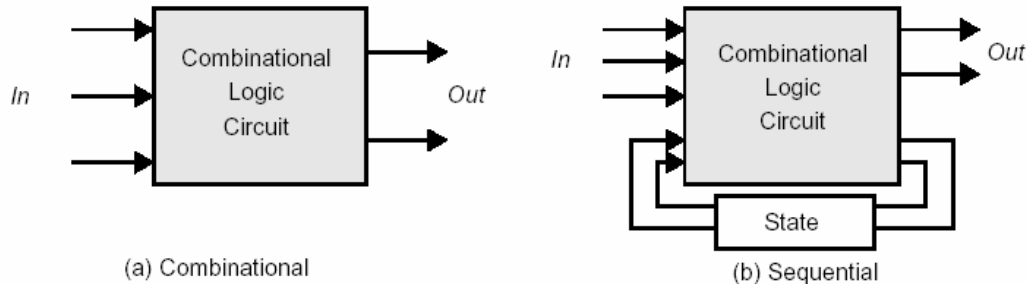
Circuit Realization

$$A \oplus B = \overline{A}B + A\overline{B} = (A + B)(\overline{A} + \overline{B}) = \overline{A\overline{B}} + \overline{\overline{A}B}$$



Combinational Logic Circuits

- Logic gates combine several logic-variable inputs to produce a logic-variable output.
- **Combinational logic circuits** are “memoryless” because their output value at a given instant depends only on the input values at that instant.



- Next time, we will study **sequential logic circuits** that possess memory because their present output value depends on previous as well as present input values.

Logical Sufficiency of NAND Gates

- If the inputs to a NAND gate are tied together, an inverter results
- From De Morgan's laws, the OR operation can be realized by inverting the input variables and combining the results in a NAND gate.
- Since the basic logic functions (AND, OR, and NOT) can be realized by using only NAND gates, **NAND gates are sufficient to realize any combinational logic function.**



Logical Sufficiency of NOR Gates

- Show how to realize the AND, OR, and NOT functions using only NOR gates

- Since the basic logic functions (AND, OR, and NOT) can be realized by using only NOR gates, **NOR gates are sufficient to realize any combinational logic function.**



Synthesis of Logic Circuits

Suppose we are given a truth table for a logic function.

Is there a method to implement the logic function using basic logic gates?

Answer: There are lots of ways, but one simple way is the “**sum of products**” implementation method:

- 1) Write the sum of products expression based on the truth table for the logic function
 - 2) Implement this expression using standard logic gates.
- We may not get the most efficient implementation this way, but we can simplify the circuit afterwards...

Logic Synthesis Example: Adder

Input			Output	
A	B	C	S_1	S_0
0	0	0	0	0
0	0	1	0	1
0	1	0	0	1
0	1	1	1	0
1	0	0	0	1
1	0	1	1	0
1	1	0	1	0
1	1	1	1	1

S_1 using sum-of-products:

- 1) Find where S_1 is 1
- 2) Write down each product of inputs which create a 1

$$\bar{A} B C \quad A \bar{B} C$$

$$A B \bar{C} \quad A B C$$

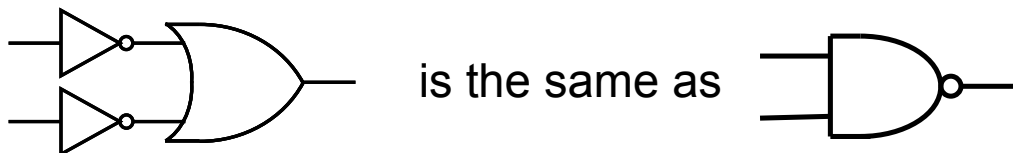
- 3) Sum all of the products

$$\bar{A} B C + A \bar{B} C + A B \bar{C} + A B C$$

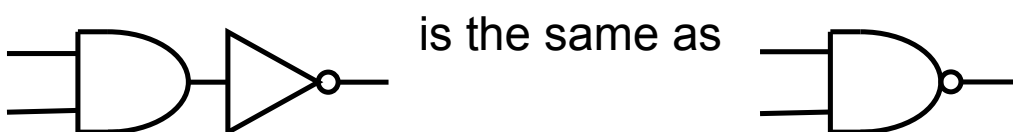
- 4) Draw the logic circuit

NAND Gate Implementation

- De Morgan's law tells us that



- By definition,



→ All sum-of-products expressions can be implemented with only NAND gates.



Creating a Better Circuit

What makes a digital circuit better?

- Fewer number of gates
- Fewer inputs on each gate
 - multi-input gates are slower
- Let's see how we can simplify the sum-of-products expression for S_1 , to make a better circuit...
 - Use the Boolean algebra relations



Karnaugh Maps

- Graphical approach to minimizing the number of terms in a logic expression:
 1. Map the truth table into a Karnaugh map (see below)
 2. For each **1**, circle the biggest block that includes that **1**
 3. Write the product that corresponds to that block.
 4. Sum all of the products

2-variable
Karnaugh Map

		B	
		0	1
A	0		
	1		

3-variable
Karnaugh Map

		BC			
		00	01	11	10
A	0				
	1				

4-variable Karnaugh Map

		CD			
		00	01	11	10
AB	00				
	01				
	11				
	10				

Karnaugh Map Example

Input			Output	
A	B	C	S ₁	S ₀
0	0	0	0	0
0	0	1	0	1
0	1	0	0	1
0	1	1	1	0
1	0	0	0	1
1	0	1	1	0
1	1	0	1	0
1	1	1	1	1

Simplification of expression for S₁:

		BC			
		00	01	11	10
A	0	0	0	1	0
	1	0	1	1	1
BC			AC	AC	AB

$$S_1 = AB + BC + AC$$

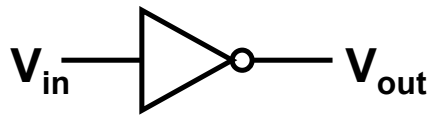
Further Comments on Karnaugh Maps

- The algebraic manipulations needed to simplify a given expression are not always obvious. Karnaugh maps make it easier to minimize the number of terms in a logic expression.
- Terminology:
 - "2-cube: 2 squares that have a common edge (-> product of 3 variables)
 - "4-cube: 4 squares with common edges (-> product of 2 variables)
- In locating cubes on a Karnaugh map, the map should be considered to fold around from top to bottom, and from left to right.
 - Squares on the right-hand side are considered to be adjacent to those on the left-hand side.
 - Squares on the top of the map are considered to be adjacent to those on the bottom.
 - Example:
The four squares in the map corners form a 4-cube

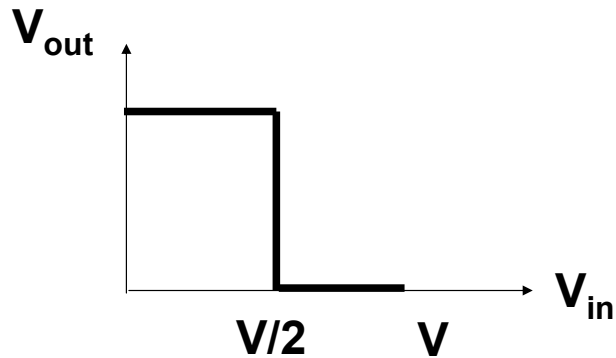
		CD			
		00	01	11	10
AB	00				
	01				
	11				
	10				



Inverter = NOT Gate

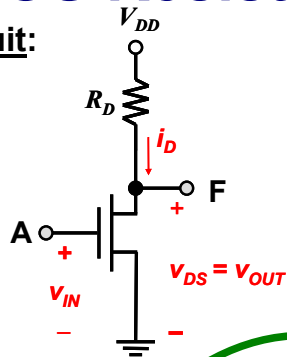


Ideal Transfer Characteristics

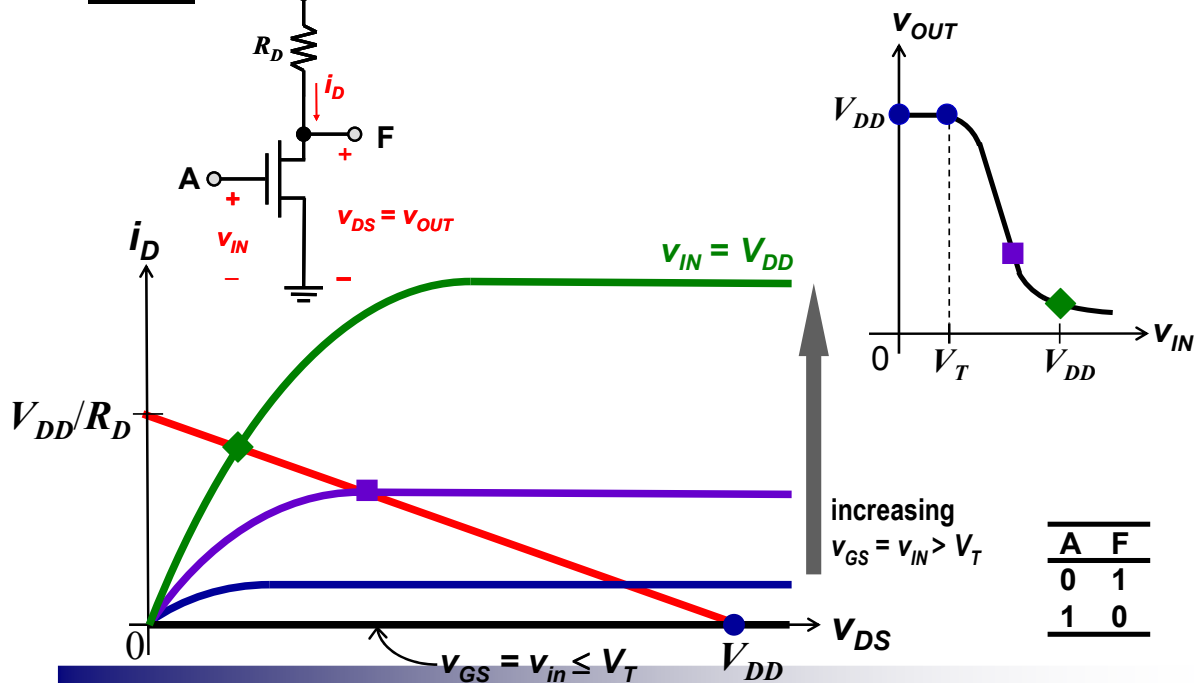


NMOS Resistor Pull-Up

Circuit:



Voltage-Transfer Characteristic

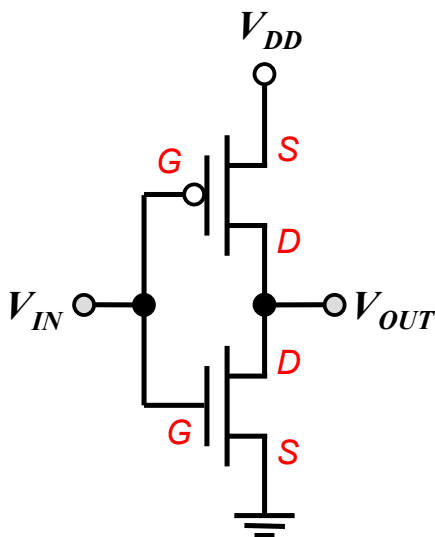


Disadvantages of NMOS Logic Gates

- Large values of R_D are required in order to
 - achieve a low value of V_{OL}
 - keep power consumption low
- Large resistors are needed, but these take up a lot of space.
- One solution is to replace the resistor with an NMOSFET that is always on.

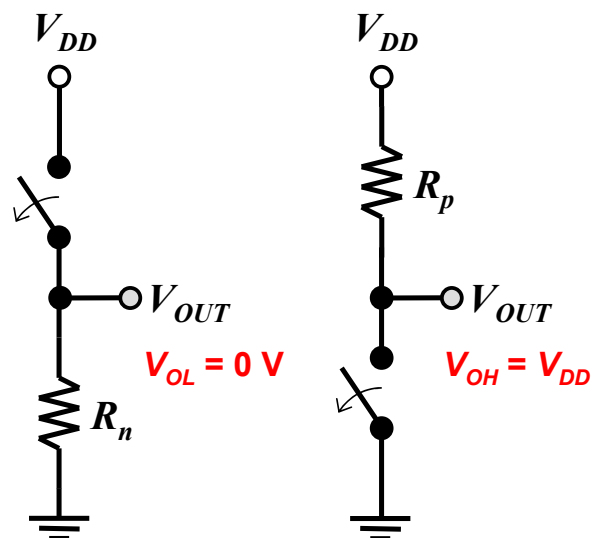
The CMOS Inverter: Intuitive Perspective

CIRCUIT



Low static power consumption, since one MOSFET is always off in steady state

SWITCH MODELS

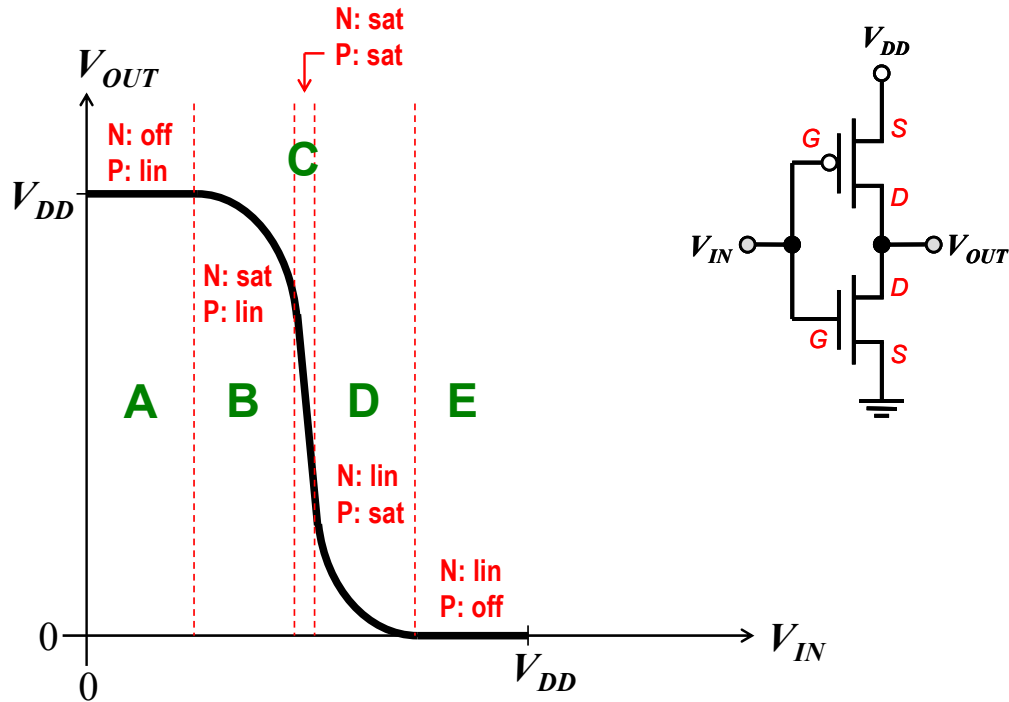


$V_{IN} = V_{DD}$

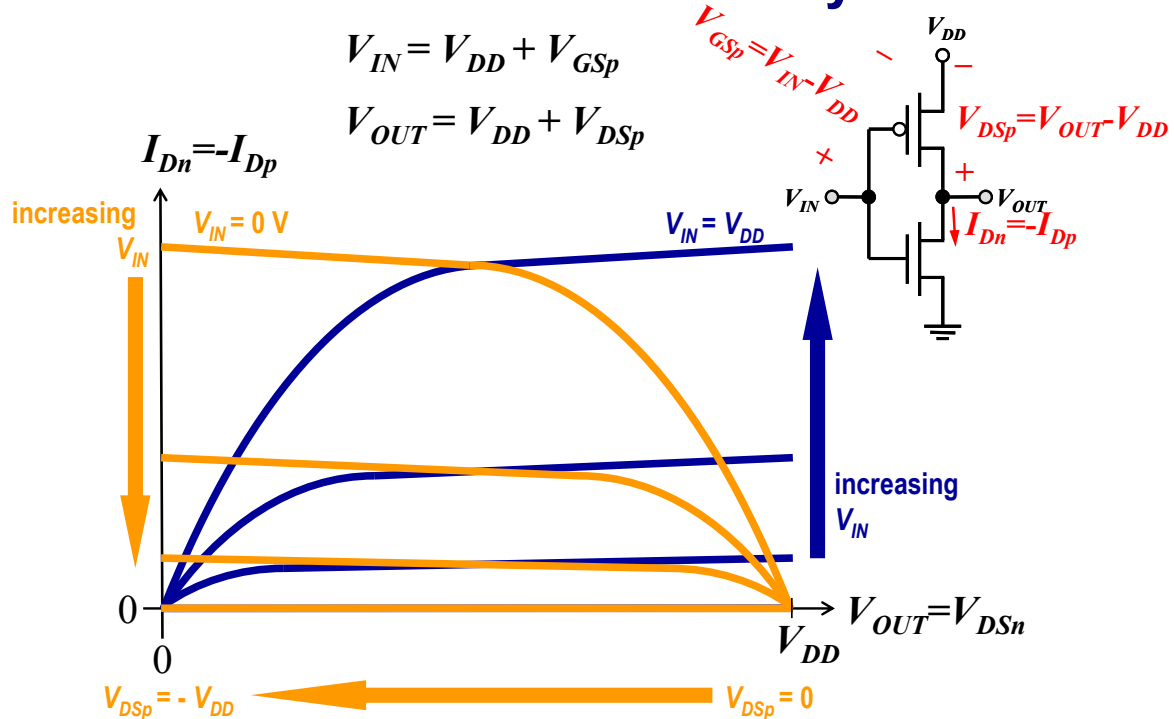
$V_{IN} = 0 \text{ V}$



CMOS Inverter Voltage Transfer Characteristic

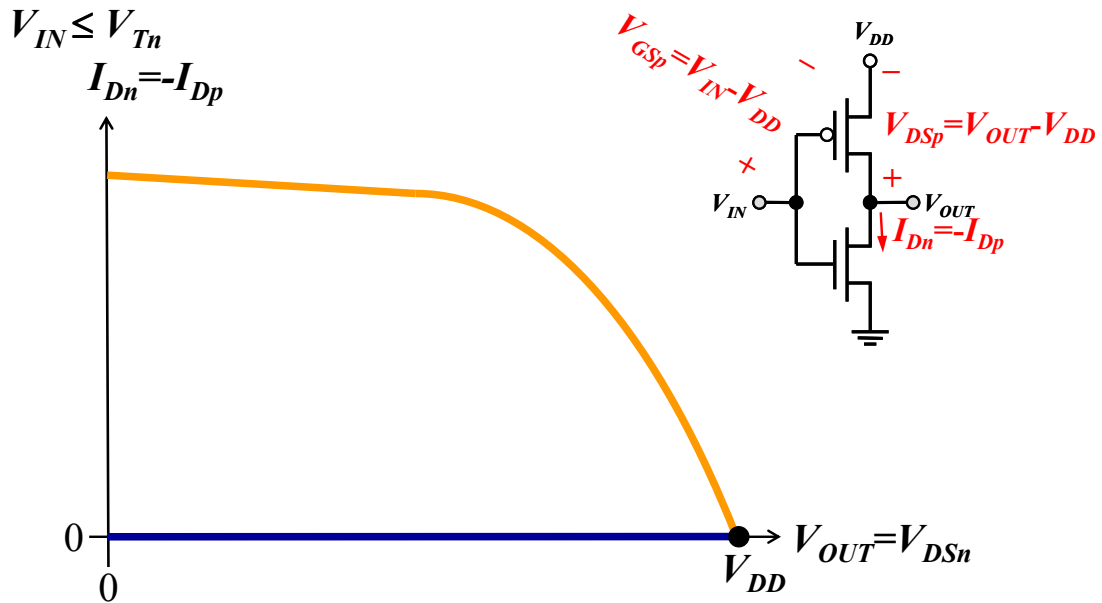


CMOS Inverter Load-Line Analysis



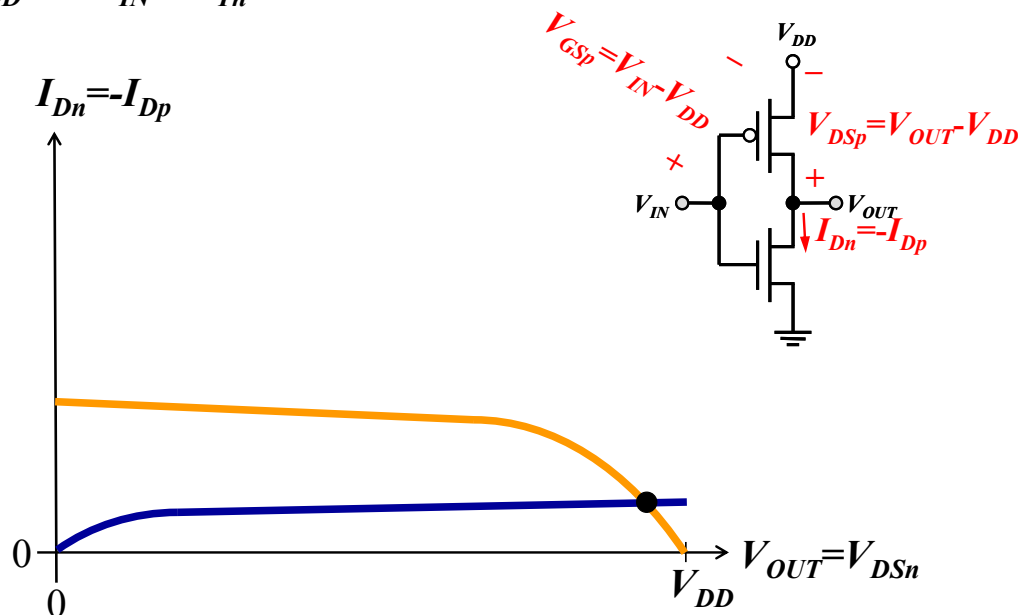


CMOS Inverter Load-Line Analysis: Region A



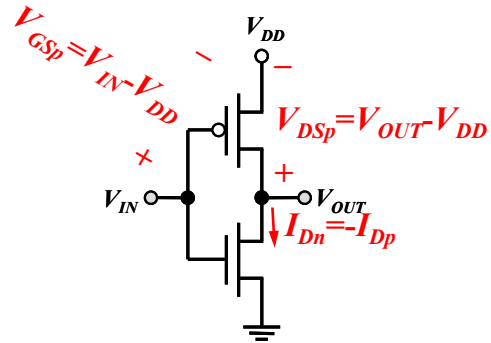
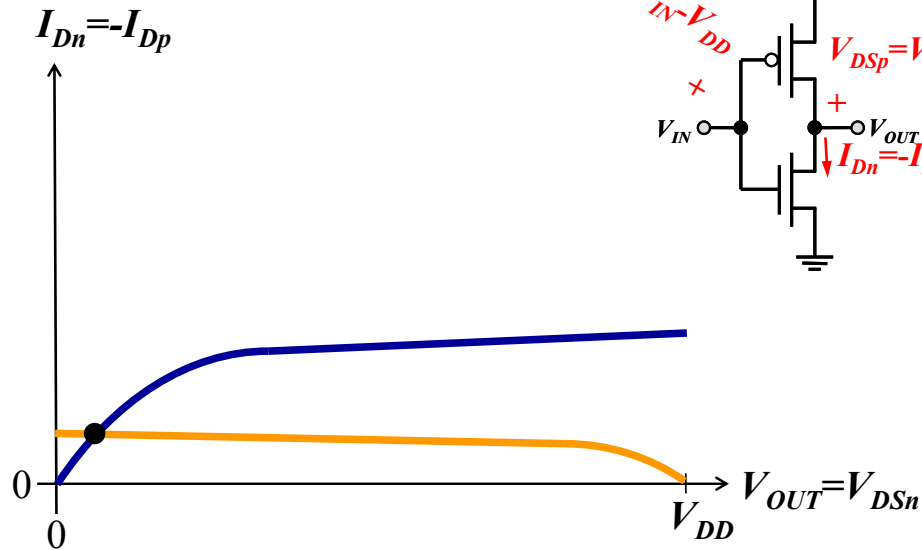
CMOS Inverter Load-Line Analysis: Region B

$$V_{DD}/2 > V_{IN} > V_{Tn}$$



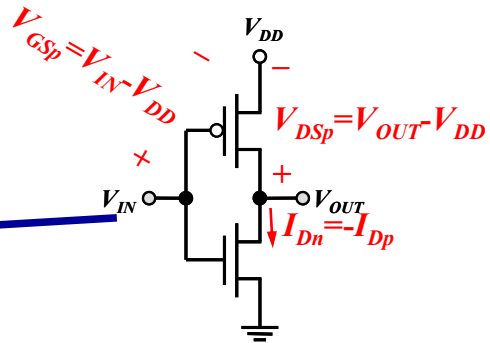
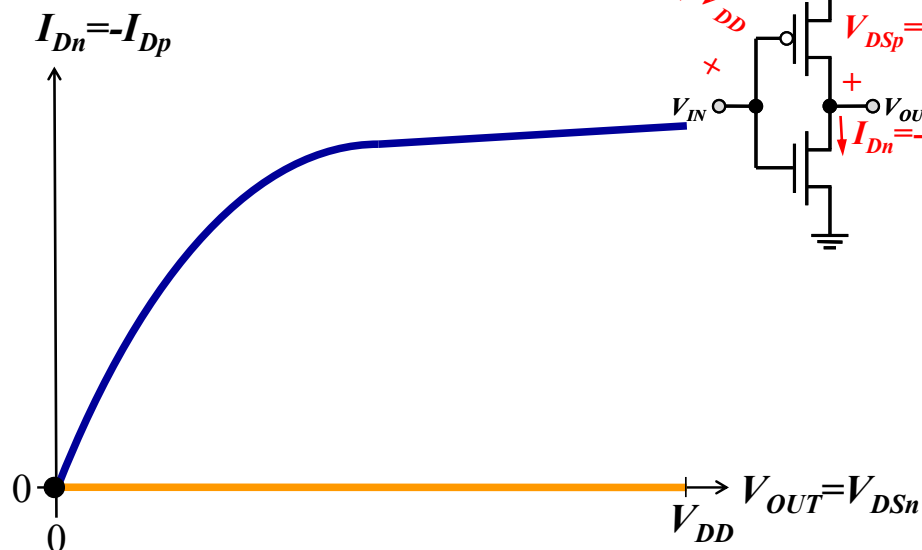
CMOS Inverter Load-Line Analysis: Region D

$$V_{DD} - |V_{Tp}| > V_{IN} > V_{DD}/2$$



CMOS Inverter Load-Line Analysis: Region E

$$V_{IN} > V_{DD} - |V_{Tp}|$$



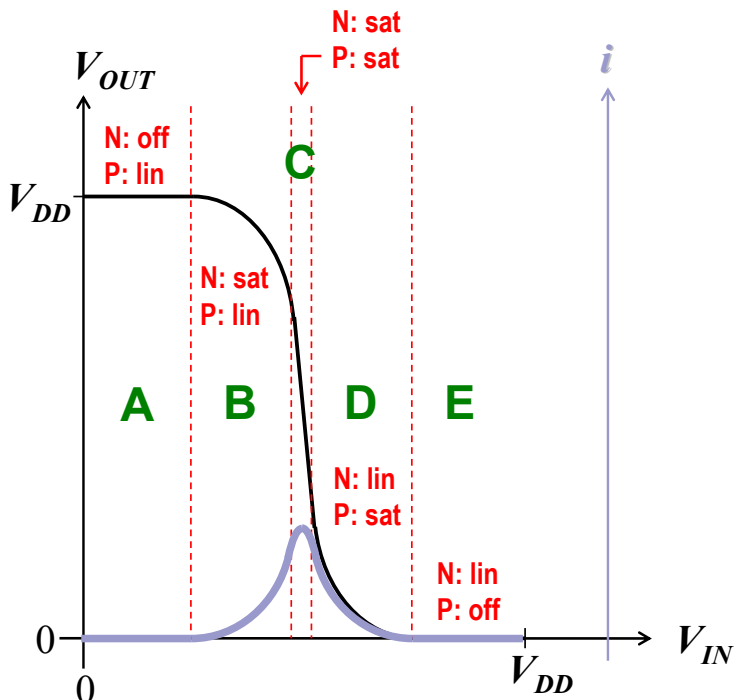
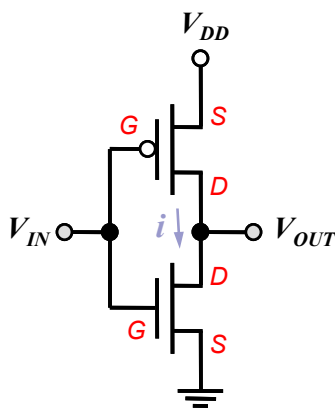


Features of CMOS Digital Circuits

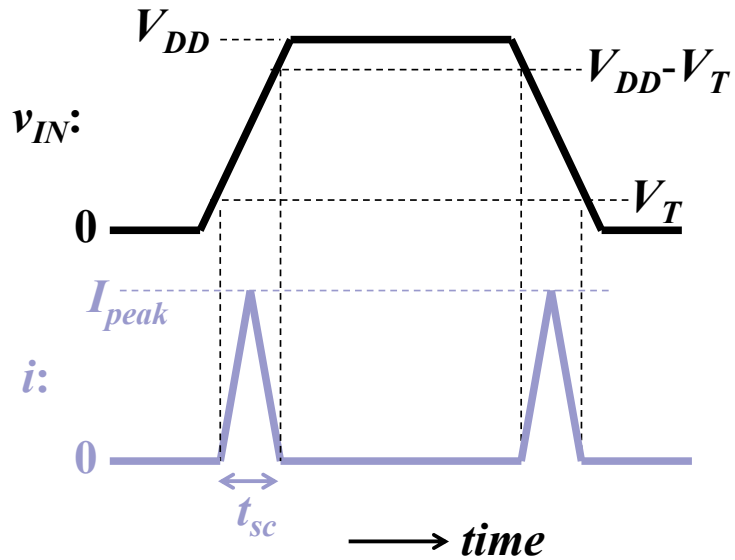
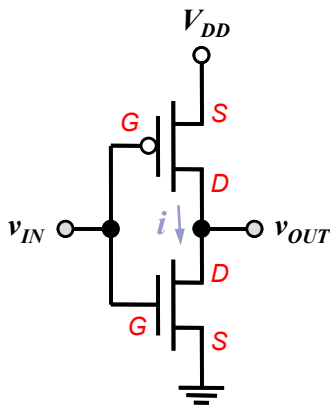
- The output is always connected to V_{DD} or **GND** in steady state
 - Full logic swing; **large noise margins**
 - Logic levels are not dependent upon the relative sizes of the devices ("**ratioless**")
- There is no direct path between V_{DD} and **GND** in steady state
 - **no static power dissipation**



The CMOS Inverter: Current Flow during Switching



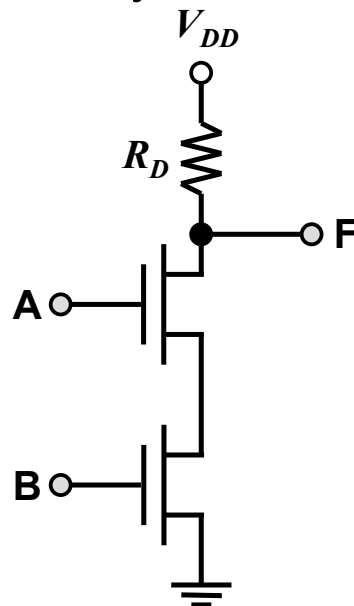
Power Dissipation due to Direct-Path Current



Energy consumed per switching period: $E_{dp} = t_{sc} V_{DD} I_{peak}$

NMOS NAND Gate

- Output is low only if both inputs are high



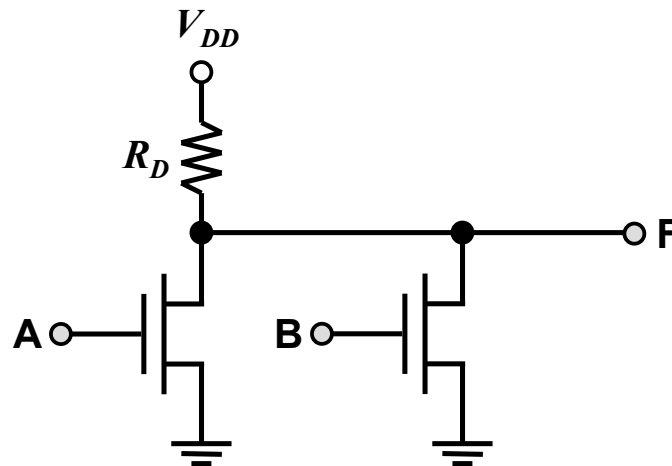
Truth Table

A	B	F
0	0	1
0	1	1
1	0	1
1	1	0



NMOS NOR Gate

- Output is low if either input is high



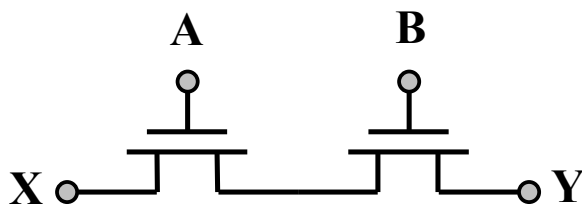
Truth Table

A	B	F
0	0	1
0	1	0
1	0	0
1	1	0

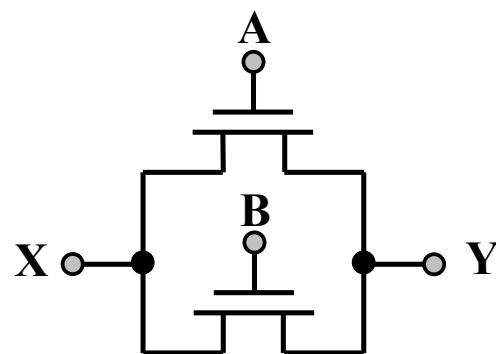


N-Channel MOSFET Operation

An NMOSFET is a closed switch when the input is high



$Y = X$ if A and B



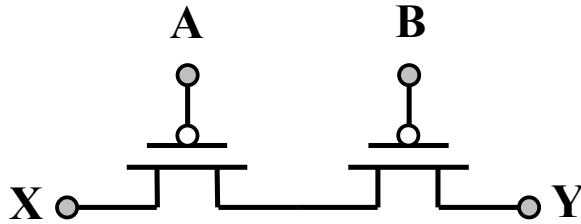
$Y = X$ if A or B

NMOSFETs pass a “strong” 0 but a “weak” 1

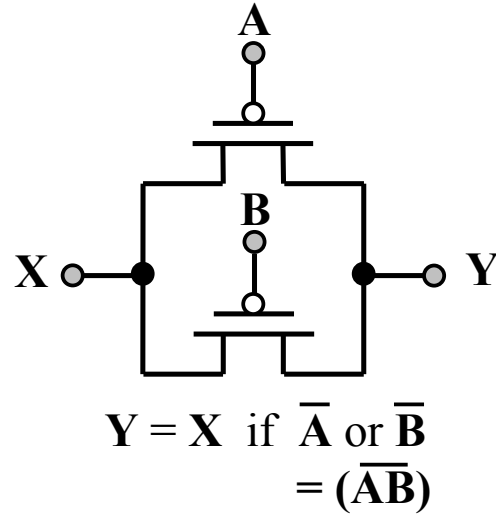


P-Channel MOSFET Operation

A PMOSFET is a closed switch when the input is low



$$Y = X \text{ if } \bar{A} \text{ and } \bar{B} \\ = (\bar{A} + \bar{B})$$



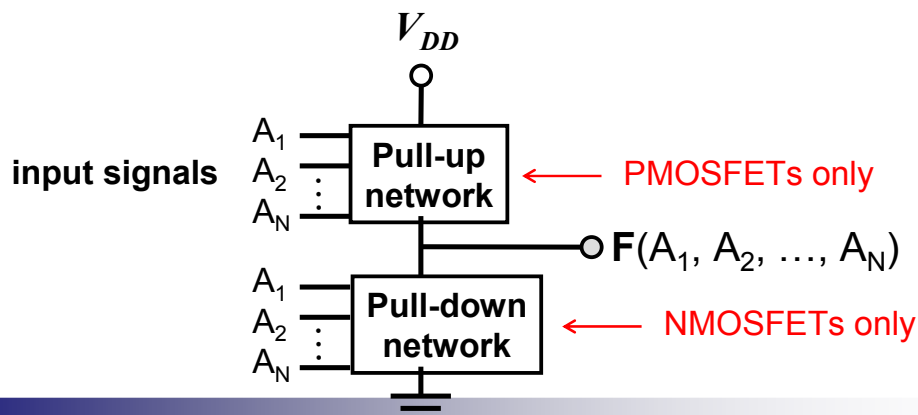
$$Y = X \text{ if } \bar{A} \text{ or } \bar{B} \\ = (\bar{A}\bar{B})$$

PMOSFETs pass a “strong” 1 but a “weak” 0



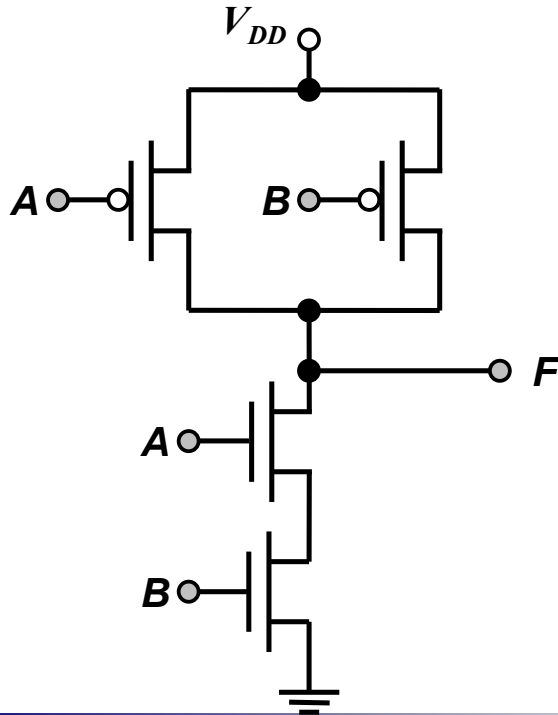
Pull-Down and Pull-Up Devices

- In CMOS logic gates, NMOSFETs are used to connect the output to **GND**, whereas PMOSFETs are used to connect the output to **V_{DD}**.
 - An NMOSFET functions as a **pull-down device** when it is turned on (gate voltage = **V_{DD}**)
 - A PMOSFET functions as a **pull-up device** when it is turned on (gate voltage = **GND**)





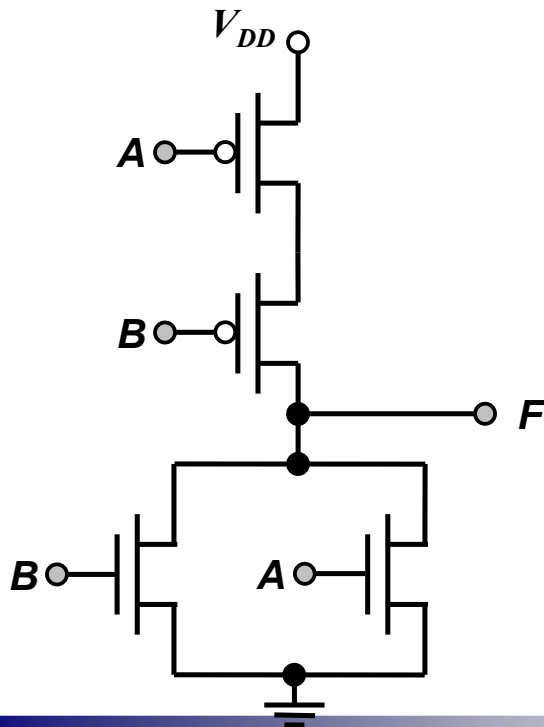
CMOS NAND Gate



A	B	F
0	0	1
0	1	1
1	0	1
1	1	0



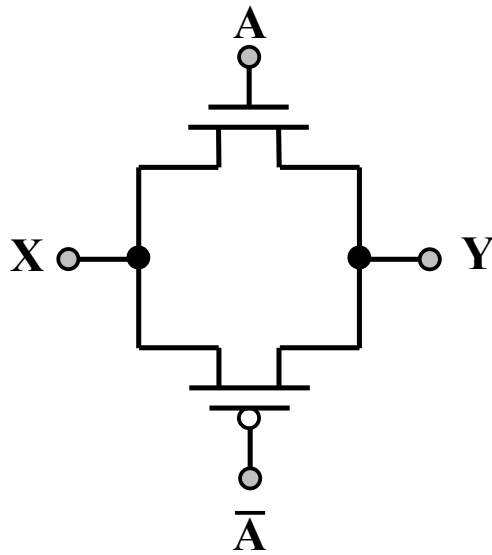
CMOS NOR Gate



A	B	F
0	0	1
0	1	0
1	0	0
1	1	0



CMOS Pass Gate



$$Y = X \text{ if } A$$



Fan in/Fan out

- Complex digital operations are formed with a variety of gates interconnected to yield the desired logic function.
- Sometimes a number of inputs are connected to one gate input and output of a gate may be connected to a number of gates.
- Fan-in: the maximum number of logic gates that can be connected at the input of a gate **without altering its performance**.
- Fan-out: the maximum number of logic gates that can be connected to the output of a gate **without altering its performance**.
- Typical fan-in and fan-out numbers are 3.